

Securing IIoT using Defence-in-Depth: Towards an End-to-End Secure Industry 4.0

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Abstract

Industry 4.0 uses a subset of the IoT, called Industrial IoT (IIoT) to achieve connectivity, interoperability and decentralisation. The deployment of industrial networks rarely considers security by design, but this becomes imperative in smart manufacturing as connectivity increases. The combination of OT and IT infrastructures in Industry 4.0 adds new security threats beyond those of traditional industrial networks. Defence-in-Depth (DiD) strategies tackle the complexity of this problem by providing multiple defence layers, each of these focusing on a particular set of threats. Additionally, the severe requirements of IIoT networks demand lightweight encryption algorithms. Nevertheless, these ciphers must provide E2E (End-to-End) security, as data pass through intermediate entities, or middleboxes, before reaching its destination. If compromised, middleboxes could expose vulnerable information to potential attackers if it is not encrypted throughout this path. With this in mind, this paper proposes a Defence-in-Depth (DiD) approach combined with the lightweight E2E encryption algorithm Attribute-Based-Encryption (ABE) and object security (i.e., OSCORE) to provide a full E2E security approach. This analysis is a critical first step to develop more complex and lightweight security frameworks suitable for Industry 4.0.

Keywords: Industry 4.0, IIoT, E2E Security, Defense in depth, OSCORE,

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1 Introduction

2 In recent years, *IoT* has become a popular term used in many areas.
3 Although there is no official definition, several attempts have been made in
4 this direction [1] [2] [3], which usually describe the IoT as a set of connected
5 devices able to process, send or receive data, with or without an Internet
6 connection. This has transformed the way people and machines communicate
7 and interact with each other. Nowadays, the IoT revolution has reached the
8 industry, leading to the fourth industrial revolution [4], or Industry 4.0.

9 *Industry 4.0* is a concept coined by the German Government [5] and pre-
10 sented in the Hannover Messe 2011. It aims to produce higher quality prod-
11 ucts and reduce production costs through the use of Industrial IoT (IIoT),
12 among other key enabling technologies. IIoT is a subset of the IoT applied
13 to industry and the evolution of industrial communications [6]. It increases
14 connectivity, interoperability and decentralisation. IIoT devices collect the
15 exchanged information en masse, which is later processed so systems can
16 carry out actions and decisions with or without human intervention. Even
17 though IoT and IIoT share some goals, their design and application envi-
18 ronment are different. For instance, the data volume that the IIoT needs to
19 manage tends to be much higher than typical IoT applications. Various re-
20 searchers have analysed the properties and constraints of IoT and IIoT [6] [7]
21 [8]. They are summarised in Table 1, where ! symbolises that it only applies
22 in particular cases—i.e., battery limitation or sleep mode, which may not
23 exist in every Industry 4.0 environment. Other features may apply to both
24 IoT and IIoT while having more relevance in the IIoT, like interdependence.
25 Uncontrolled alterations in actuators, sensors and control systems may risk
26 the availability of the entire system. Interdependence is not as critical in the
27 IoT, where nodes join and leave networks often. Such aspects must be con-
28 sidered during industrial systems design phase since they cause a significant
29 impact on security and communications, as do battery and computational
30 limitations. Note that these features are so restrictive that they have the
31 potential to condition the entire network, even if they only affect a few nodes
32 in the network.

33 Because of the constrained nature of IIoT devices, sometimes data pro-
34 cessing is carried out in edge devices or the Cloud [9]. Thus, wireless commu-

	IoT	IIoT
Battery Limitation	✓	!
Computing Limitation	✓	✓
Sleep-Mode	✓	!
Interdependance	✓	✓
Heterogeneity	✓	✓
Structured Nodes	×	✓
Scalability	✓	✓
Interoperability	✓	✓
Very High Data Volume	×	✓

Table 1: Feature comparison between IoT and IIoT.

35 nications are increasingly common in industrial environments, using proto-
 36 cols such as Zigbee, WirelessHART, Trusted Wireless, WiFi or Bluetooth
 37 [10]. The application-layer protocols running on top of them should be
 38 lightweight and address the constrained nature of IIoT devices. Therefore,
 39 protocols typically designed for IP networks may not be suitable for the IIoT.
 40 In this context, IETF Working Group, CoRE [11], has proposed a framework
 41 for applications that run on constrained devices and networks. The lightness
 42 of their solution might be of particular interest in smart manufacturing, where
 43 where IIoT devices exchange substantial volumes of information.

44 Industry 4.0 architectures are decentralised systems, in which messages
 45 go through proxies, gateways and other middleboxes to save bandwidth and
 46 memory or perform protocol-translation operations [12]. These middleboxes
 47 provide scalability, efficiency and interoperability among nodes. However,
 48 they have full access to the relayed data, even if communications have been
 49 protected with transport-layer security (TLS). This might cause security in-
 50 cidents if they are compromised, in which case TLS is not enough. Instead,
 51 additional end-to-end (E2E) security mechanisms, capable of guaranteeing
 52 that data is not exposed to third parties, are required. Additionally, due to
 53 the long life span of the Operational Technology (OT) devices, legacy related
 54 issues must be considered. Otherwise, the limitations of these devices might
 55 cause various incidents, e.g., safety violations, monetary losses or information
 56 theft.

57 With this in mind, the purpose of this paper is to study the existing se-
 58 curity measures for Industry 4.0 and explore options to ensure E2E security
 59 in such environments. Then, we propose a secure Industry 4.0 framework

60 that provides E2E security combining Defence in Depth (DiD) techniques,
 61 application-layer security and functional encryption. These concepts are ex-
 62 tensively explained throughout the paper.

63 The remaining paper is structured as follows: Section 2 presents an
 64 overview industrial security, points out the most relevant Industry 4.0 se-
 65 curity requirements and provides security best practices for such scenarios.
 66 Section 3 introduces the goals of any DiD strategy and proposes DiD lay-
 67 ers compliant with them, as well as an example of a network segmentation
 68 scheme. Section 4 analyses the need and implications of using encryption in
 69 manufacturing, and how it can be used to obtain E2E security. Section 5
 70 and Section 6 introduce object security (i.e., OSCORE) and ABE and discuss
 71 their applicability in Industry 4.0 scenarios. Finally, Section 7 highlights the
 72 most important insights and concludes the paper.

73 **2. Security in Industry 4.0: A general approach**

74 Industry 4.0 uses other enabling technologies that go beyond IIoT. In the
 75 case of manufacturing, systems are complex structures formed by Information
 76 Technology (IT) and Operational Technology (OT) networks. IT networks
 77 refer to the technologies used for information processing and telecommu-
 78 nications equipment. OT networks are related to industrial equipment re-
 79 sponsible for monitoring and controlling physical devices. Effective security
 80 architectures should be included since the system design stage and reviewed
 81 often [13]. They should also take into account the growing connectivity of
 82 OT networks, which makes them resemble IT networks more than ever, while
 83 still needing to remain separated, e.g., by keeping IT and OT infrastructures
 84 separate using New Generation Firewalls (NGFWs). These Firewalls offer
 85 features like application-level inspection and a designated update path, which
 86 enhance network security and ease security updates. In terms of security, OT
 87 and IT have different priorities, as seen in Table 2.

Priority Level	OT	IT
1	Availability	Confidentiality
2	Integrity	Integrity
3	Confidentiality	Availability

Table 2: Prioritisation of security requirements for IT and OT networks.

88 Differences between OT and IT have been widely studied in the literature

89 and are not the focus of this paper. Still, addressing them is important to
 90 understand why traditional IT security approaches cannot be directly applied
 91 to OT networks. Their most relevant traits from a security point of view are
 92 shown in Table 3, which summarises the analysis presented in [14]. It is of
 93 particular relevance to highlight the strict latency requirements, the need for
 94 a fault-tolerant design or the much longer lifetime of OT systems compared
 95 to IT systems. These particularities should be considered when adapting
 96 existing IT solutions to the OT environment. For instance, Defence in Depth
 97 (DiD) strategies.

	OT	IT
Performance requirements	Real-Time Delays unacceptable	No Real-Time Delays acceptable
Fault-Tolerance	Essential	Not important
Updates	Should first be implemented in a controlled environment	Updates are straightforward
Communications	Proprietary protocols Wired and Wireless Complex Networks	Standard protocols Wired networks IT networking practices
Lifetime	10-15 years	3-5 years
Device Location	May be remote and isolated	Local and easy to access

Table 3: Summary of OT and IT networks differences [14].

98 2.1. General Security Recommendations

99 Unfortunately, poor security practices have been discovered in industrial
 100 networks, like those emulated in [15]. These security flaws particularly affect
 101 small business without IT staff, which do not have the required knowledge or
 102 resources to invest in strong security mechanisms and equipment. However,
 103 it is important to follow at least the next recommendations:

- 104 • Keep software up-to-date: Enterprises sometimes use hardware with
 105 known vulnerabilities, e.g., Allen-Bradley’s MicroLogix [16] [17] or Sie-
 106 mens Simatic [18]. To patch them, it is recommended to apply the
 107 security updates provided by the original manufacturers as soon as
 108 they are made available. To minimise the effects on production, up-
 109 dates should be applied first in a controlled environment simulating the
 110 real one. However, occasionally manufacturers may refuse to offer an
 111 update if the vulnerable device has reached the end of its life-cycle. In
 112 that case, other approaches, such as hardening, might be studied.

- 113 • Use strong passwords: Passwords for HMIs (Human-Machine Inter-
114 faces) and workstations should be complex and unique, and they should
115 never be the default ones. VNC (Virtual Network Computing) systems
116 should have specific passwords for remote control. Basic recommenda-
117 tions for them is having a minimum of 8 characters, with a combination
118 of capital and lower cases, special characters and numbers. Under no
119 circumstances should these passwords be related to the identity of the
120 device they protect.
- 121 • Implement strict access control mechanisms: Having some kind of ac-
122 cess control for the mentioned HMIs and workstations is strongly rec-
123 ommended. A similar approach should be considered when dealing
124 with file servers.
- 125 • Implement network segmentation: Unrelated networks should have
126 physical and logical separations. This is extensively explained in Sec-
127 tion 3.2.

128 Following these recommendations enhances security by decreasing some
129 of the most well-known vulnerabilities. However, most industrial systems
130 require more complex security measures, which will be used to fulfil the
131 security requirements defined in the next Section.

132 *2.2. Industry 4.0 Specific Security Recommendations*

133 The particularities of industrial manufacturing add additional constraints
134 in the design of efficient security approaches for OT networks. Nevertheless,
135 the traditional security requirements of IT should still be guaranteed in in-
136 dustrial security. They are authentication, confidentiality, access control,
137 integrity, non-repudiation and availability. The following recommendations
138 address each of them:

- 139 • Availability: To guarantee this requirement, the system should be de-
140 signed with fault-tolerance in mind. Critical devices and networks
141 should have a redundant counterpart to replace the original in the
142 event of failure or security breach. These redundancy mechanisms help
143 prevent DoS (Denial of Service) attacks and assure users' safety.
- 144 • Authentication and authorisation: According to the IEC 62443-4-2 [19],
145 every user in a system has to be authenticated, and every requester of

146 an operation needs to be previously authorised. The advised way to
147 achieve this [14] is with the use of whitelists and only allow communi-
148 cations between authenticated and authorised source-destination pairs.

- 149 • Access control: This must be considered when accessing devices' con-
150 figuration and any resource in the network. Role-based access controls
151 are strongly recommended [14]. The aim is to diminish the effects of
152 impersonation attacks and favour confidentiality. This is of especial
153 relevance in control systems and databases. Preventing attackers from
154 accessing databases also prevents them from getting critical informa-
155 tion and credentials that could later be used to access critical control
156 systems.
- 157 • Integrity and confidentiality: Unwanted message modification can have
158 dangerous consequences for systems and users in the IIoT. For instance,
159 as [20] presents, exposing or maliciously modifying sensitive informa-
160 tion may put a persons' life in danger in case of a health emergency.
161 Thus, data has to remain unchanged and confidential during capture,
162 retrieval, update, storage and transport. Only authorised users should
163 be able to read or modify it. For example, as shown in Section 6, by
164 using ABE only users with specific attributes or roles would be able to
165 access the encrypted information.
- 166 • Non-Repudiation: This guarantees that messages are transmitted in a
167 way that the authenticity of the information cannot be questioned later
168 [21]. It is especially relevant in Human User Interfaces [19], so human
169 actions are reflected in the system and can be traced back to the user.

170 Besides implementing the above-mentioned security measures, a layered
171 security approach is strongly encouraged. In the coming section, we introduce
172 the concept of Defence in Depth (DiD) applied to Industry 4.0 infrastruc-
173 tures.

174 **3. Security in Industry 4.0: A DiD approach**

175 One of the advanced techniques to secure industrial environments is De-
176 fence in Depth (DiD). According to the IEC 62443-4-1 [22], the goal of this
177 approach is to limit the damage in case of an attack by implementing lay-
178 ered security controls. DiD is an effective security method that addresses

179 many attack vectors, as each layer provides additional defence mechanisms.
180 It can be implemented in both OT and IT networks with different security
181 techniques but similar goals.

182 *3.1. DiD Goals*

183 Most enterprises are familiar with IT security, but not so much with OT
184 security. Until recently, the only access points to the systems were physical
185 and security was not a concern. With the evolution of the industry to Indus-
186 try 4.0 and the growing connectivity of the systems, cybersecurity becomes
187 a requirement to be implemented as part of the systems' design. Various
188 institutions worldwide such as the NIST [14], the Spanish INCIBE [23], and
189 even standards as the IEC 62443-4-1 [22] and IEC 62443-4-2 [19] have ad-
190 dressed the topic of security. As [13] points out, this may cause a flood of
191 information about how to integrate them in different organisations. Still,
192 these guidelines and standards have some common points, and from them,
193 the desired goals for a DiD strategy can be drawn. Regardless of which layers
194 are implemented in the DiD strategy, they should always meet the following
195 objectives and procedures:

- 196 • The security requirements of Section 2.2. Availability is the main prior-
197 ity. Regarding data integrity, it can be compromised accidentally or as
198 a result of an attack. The first case can be the result of interferences in
199 industrial communications and measures to guarantee integrity are al-
200 ready used (i.e., CRC). However, these measures may not be enough to
201 handle active attacks, which may result in sabotage. Instead, a combi-
202 nation of role-based access control, encryption and integrity preserving
203 algorithms (i.e., digital signatures) should be used.

- 204 • Restricted physical and logical access to the system, taking into ac-
205 count both external and internal threats. The connection between OT
206 and IT should be restricted, and following the recommendations of
207 [14], achieved using a demilitarised zone (DMZ) and reducing traffic to
208 specific and documented services and ports. The use of DMZs in com-
209 bination with unidirectional gateways and firewalls restrict the logical
210 access to the system and help achieve the restricted data flow required
211 in the IEC 62443-4-2 [19]. To restrict physical access, it is advised
212 to use Biometric Systems and Smart Cards. The access permissions
213 should be implemented following a least-privilege approach and issued

214 by a trusted entity. This entity should also keep them up-to-date, to
215 reflect the current situation and prevent security breaches.

216 • ICSs protection from known vulnerabilities. The long lifetime of these
217 devices makes them particularly vulnerable to attacks. Updates and
218 security patches should be installed as explained in Section 2.1. In
219 case no more security upgrades are available, a vulnerability assess-
220 ment should be performed and a rigorous hardening process should be
221 considered, e.g., using whitelists, reducing application services to the
222 minimum or restricting users' privileges and roles as much as possible.

223 • System monitoring and security incidents detection. Malfunctioning
224 ICS and misconfigured services create vulnerabilities in the systems.
225 Detecting them on time can prevent future security attacks. The im-
226 plementation of Intrusion Detection Systems (IDS) or Intrusion Protec-
227 tion Systems (IPS) is recommended to detect possible threats as soon
228 as possible. These systems detect abnormal behaviours by comparing
229 the current and expected status. This way attackers can be blocked
230 while attempting to enter the system.

231 • Periodical evaluations of security. Following the guidelines of [14], se-
232 curity should be addressed during the design, use, maintenance and
233 removal of industrial systems. This includes hardware, software and
234 security policies.

235 • Limit the impact on production. Essential functions that guarantee
236 health, safety, environment maintenance and equipment availability
237 [19] cannot be negatively affected by security measures or emergen-
238 cies. Therefore, it is essential to find a balance that gives the system
239 as much security as possible, while still fulfilling all the production
240 requirements. Besides, since not every attack can be prevented, fast
241 restoration plans are recommended to be in place too.

242 • Isolation of critical systems. ICSs and control networks should have
243 no connection to the Internet, not even through firewalls. However, in
244 case this is strictly necessary, communications must use only proved
245 secure protocols and go through a DMZ.

246 Achieving these goals can be eased when applied in combination with
247 network segmentation, first mentioned in Section 2.1. It is required by IEC

248 62443-4-2 [19] and increases security by separating the network both logically
249 and physically.

250 *3.2. Proposed DiD Layers*

251 Network segmentation enhances availability [14] and improves the sys-
252 tem's reliability [19]. Segmentation can both be physical or logical (e.g.,
253 gateways, firewalls, VPNs, VLANs), which might be implemented from the
254 link-layer up to the application layer. Logical segmentation is more flexible
255 and easier to implement but it may be bypassed and lead to single-points-
256 of-failure, while physical segmentation is more secure but also more complex
257 and expensive [19]. Thus, segmentation techniques should be analysed on a
258 case-per-case basis since there is no universal solution.

259 The key to successful security frameworks lies in the combination of net-
260 work segmentation (Figure 1) and a DiD approach. Each of the security
261 zones should consist of assets with similar security needs, thereby facilitat-
262 ing monitoring and logical access control. The zones can also be subdivided
263 into more segments as needed, improving overall security. In agreement with
264 the IEC 62443-4-1 [22], the DiD layers should provide additional defence
265 mechanisms by supporting the secure design principles specified in the same
266 standard. The choice of which mechanisms to implement in each layer is
267 left to the user-e.g., IDSs, IPSs, firewalls, security gateways or encryption al-
268 gorithms. Thus, following those guidelines along with the required network
269 segmentation of the IEC 62443-4-2 [19], a DiD layered approach is presented
270 in Figure 2, where each layer has the following purposes:

271 *3.2.1. Physical Security*

272 The first security layer handles physical security. Measures to ensure re-
273 stricted physical access must adapt to the particularities of the organisation.
274 As introduced in Sec 3.1 smart cards and biometric systems are potential
275 solutions. It should be taken into account that although Figure 2 presents
276 physical security as a single layer, this security layer is distributed through-
277 out the enterprise infrastructure, and therefore it may include a wide variety
278 of security mechanisms. Context-dependant access may be necessary. For
279 instance, access to locations like the control room or the general assembly
280 line may vary depending on the hour or user-role. Physical security is of
281 crucial importance since this is the first layer of protection against external
282 attacks.

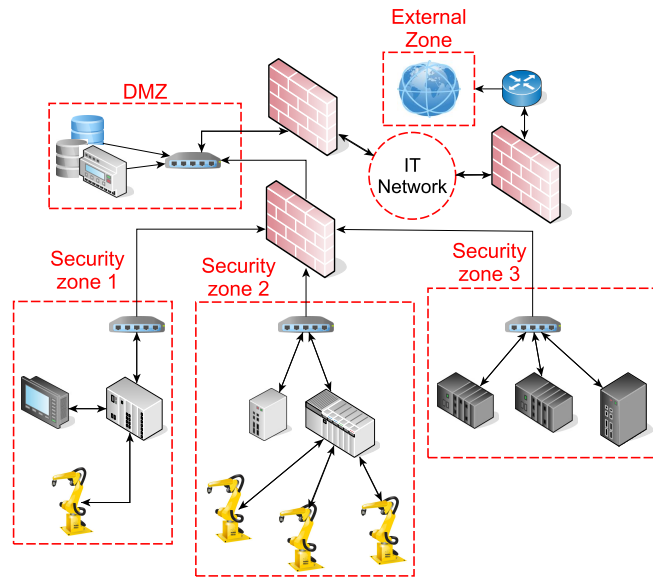


Figure 1: OT network segmentation with three security zones and a DMZ separated by firewalls.

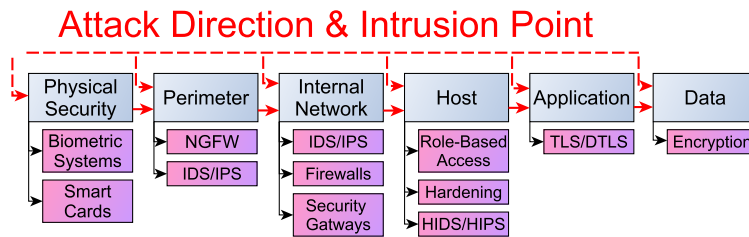


Figure 2: Security layers in DiD (in blue) with their corresponding security measures (in pink).

283 3.2.2. Perimeter

284 Perimetral security is the layer that protects the OT network from ex-
 285 ternal communications by restricting the access and filtering unauthorised
 286 communications, including the ones coming from the IT network. A com-
 287 mon way of achieving this has been limiting traffic to specific ports. However,
 288 smart manufacturing needs to manage a much higher volume of traffic, while
 289 the equipment may still be old. Thus, it is possible to flood a legacy system
 290 by accident and cause a DoS attack. To prevent this, solutions based on Next
 291 Generation Firewalls (NGFWs) should be implemented. These firewalls can

292 be used as shown in Figure 1. In it, the IT and OT networks are separated
293 by a DMZ that will filter every communication between both networks, and
294 which is placed between two NGFW. These firewalls, as mentioned in Sec-
295 tion 2, offer deep-packet inspection and IDS/IPS functionalities, becoming
296 very useful for network monitoring and traffic filtering tasks. Filtering is
297 recommended to be performed following a whitelisting approach. Although
298 whitelisting may not be feasible in every firewall, it must be used in high-
299 risk security environments. Meanwhile, monitoring can be active or passive,
300 depending on the particular requirements of the system. If the purpose is
301 to analyse incidents and learn about attack patterns to evolve the security
302 infrastructure, IDS would be sufficient. Instead, if the aim is to stop the
303 intrusion as soon as possible without any further analysis, IPS ought to be
304 used. It is important to note that applying an IPS approach requires a deep
305 knowledge of the network traffic, since an IPS reacting to a false positive
306 may lead to an unexpected DoS. Note also that firewalls and IDS systems
307 are complementary technologies, and one does not substitute the other.

308 *3.2.3. Internal Network*

309 So far, the proposed security layers protect the system as a whole and are
310 designed to avoid unauthorised network accesses from the outside. In con-
311 trast, the following security layers are devised to protect network resources
312 when attackers are already within the network. Thus, they will be applied
313 independently to security zone or sub-network. Because stopping sophisti-
314 cated attacks requires more complex security measures, applying them to
315 smaller networks improves their efficiency and allows them to be specifically
316 designed with the sub-network requirements in mind. This level of protec-
317 tion is mainly composed of devices that control the sub-network inbound and
318 outbound traffic, such as IDSs/IPSSs, firewalls and security gateways.

319 *3.2.4. Host*

320 The goal of the next layer is to protect each of the devices inside a security
321 zone. This is of particular relevance in OT security, where targeted attacks
322 on critical systems may cause significant damage to the whole system. Thus,
323 it is crucial to detect anomalies by actively scanning for vulnerabilities and
324 modifications in the firmware or device configuration. The security measures
325 applied in this layer vary depending on the system's capabilities and limi-
326 tations. If newer devices support role-based access control, it is advisable
327 to apply it. This measure can be reinforced by following the recommended

328 practices in Section 2.1 and the hardening practices introduced in Section
 329 3.1. In case the system cannot implement advanced authentication mecha-
 330 nisms, reinforced access control should be considered. If the asset supports
 331 them, additional security measures at host level can also be considered, such
 332 as host-based IDS (HIDS) or host-based IPS (HIPS). These would provide
 333 another layer for monitoring and detection of abnormal situations in the host.

334 3.2.5. Application and Data

335 These layers are the last safeguards against attacks, and the most related
 336 to IT security. They aim to protect data and services from attacks that
 337 have not been detected by the previous layers. It is strongly recommended
 338 to use strong application-level security mechanisms whenever possible, along
 339 with data encryption. Even if they remain independent, these layers are
 340 closely related, as the encryption protocol choice may be determined by the
 341 application protocol. This will be further explained in Section 4. Application
 342 and data layers should also deal with remote accesses, which ought to be
 343 controlled. This can be done with secured VPNs, a temporal user in secured
 344 PCs or by subjecting accessing users to vulnerability scans.

345 The proposed DiD layers fulfil the requirements of Section 3.1, as shown
 346 in Table 4, and accomplish all the goals of a DiD strategy, some even in
 347 more than one layer. Despite this redundancy, the IEC 62443-4-1 DiD rec-
 348 ommendations are fulfilled since the layers remain autonomous and similar
 349 functionalities are achieved by different means. Thus, if an attacker breaks
 350 into the system, they still have to surpass many security barriers with differ-
 351 ent weaknesses before achieving their goal.

		Physical Layer	Perimeter	Internal Network	Host	Application	Data
Restricting Physical Access		●	○	○	○	○	○
Restricting logical access	<i>To Network</i>	○	●	●	○	○	○
	<i>To Devices</i>	○	○	○	●	○	○
Hardening		○	○	○	●	○	○
Protecting unwanted modification of data	<i>Role-Based Access</i>	●	○	○	●	○	●
	<i>Encryption</i>	○	○	○	○	●	●
Monitoring		○	●	●	●	○	○

Table 4: Goals covered by the proposed security layers. ○No ; ●Yes; ●Some cases

352 In summary, Industry 4.0 requires that IT and OT work together from the
353 design stage on behalf of network security. For this purpose, passive mecha-
354 nisms such as access control, traffic analysis and intrusion detection should be
355 combined with active mechanisms like traffic filtering, vulnerability scanning
356 and hardening. It is also of the utmost importance to provide the informa-
357 tion collected throughout all these layers, clearly and comprehensively, to
358 deal with potential problems as soon as possible. Finally, all of these mech-
359 anisms must be applied with consideration of network segmentation. Every
360 middlebox or node used to connect assets and capable of communication
361 is likely to have full access to data, so E2E security measures ought to be
362 studied and implemented.

363 **4. Encryption for Industry 4.0**

364 Industry 4.0 deals with a lot of sensitive information related to the man-
365 ufacturing process and the workers involved in it. Therefore, maintaining
366 data confidentiality is vital to any Industry 4.0 security architecture, and it
367 is achieved with cryptography. However, IIoT devices (e.g., smart robots,
368 gateways, sensors or actuators) are heterogeneous in terms of memory, com-
369 munication and processing capabilities. These constraints must be taken into
370 account since encryption and decryption are computationally expensive op-
371 erations and may introduce latencies. Lightweight encryption ciphers, origi-
372 nally devised for the IoT, may be suitable for the IIoT. As was introduced in
373 [13] IoT security techniques may be applied to smart manufacturing, as long
374 as the particularities of the new domain are addressed. Thus, although there
375 are challenges to applying encryption in industry, there are also mechanisms
376 to reduce its impact as long as network security requirements and computing
377 limitations are taken into account. For instance, asymmetric cryptography
378 requires a high amount of computing and memory resources compared to
379 symmetric cryptography, and it is best suited for administrative purposes
380 [14]. Meanwhile, symmetric cryptography can be applied to the data stream
381 and network traffic [14], but it involves sharing a key beforehand, and this
382 is not always possible [6]. Finally, it is important to note that not every
383 IIoT node has encryption capabilities. While some are able to perform state-
384 of-the-art encryption, others may not have the processing power for it. In
385 this case, relegating cryptography to hardware accelerators [14] may be the
386 only available solution. In any case, encryption is encouraged to be included
387 in the design of E2E security architectures whenever possible, especially in

388 wireless networks.

389 4.1. Towards E2E Security

390 Section 3.2 shows the need to introduce intermediate entities (like gate-
391 ways and proxies) to achieve security in network segmentation. IIoT devices
392 may use lightweight communication protocols, such as MQTT [24] or AMQP
393 [25], and these need to be translated to protocols specially designed for indus-
394 trial purposes (e.g., Profibus, Profinet, Ethernet/IP or EtherCAT). Protocol
395 translation takes place in gateways that need access to the data, so messages
396 have to be constantly decrypted and encrypted again. Therefore, communi-
397 cation security is broken at every middlebox (Figure 3) and instead of E2E
398 security (i.e., secure communication is guaranteed from the sender to the final
399 destination, Figure 4), there is hop-by-hop security, which does not maintain
400 the required confidentiality if the intermediate entities are compromised.

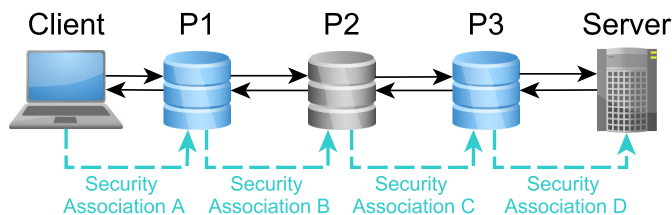


Figure 3: Hop-By-Hop Security. Security is guaranteed for every security association, but not from Client to Server.

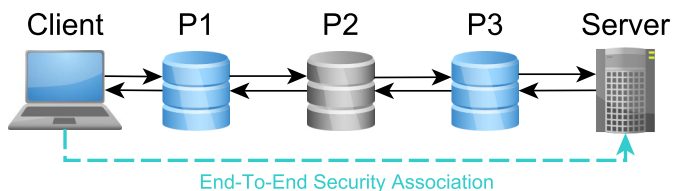


Figure 4: E2E Security. Middleboxes only have access to the information they need to forward the message to the next endpoint.

401 E2E security requires maintaining confidentiality and integrity up to the
402 destination while allowing proxies and gateways to do their jobs. For this
403 to happen, these devices should only have access to the indispensable parts
404 of the message, while the rest is hidden from them. Typically, asymmetric

405 and symmetric encryption schemes view encryption as an all-or-nothing op-
406 eration (i.e., the user either decrypts the entire message or learns nothing
407 about it [26]). Thus, middleboxes would get too much information, making
408 these ciphers not the best suited for decentralised architectures. As such,
409 it might be necessary to encrypt data so it can be shared at a fine-grained
410 level. This can be achieved with object security [27], which would encrypt
411 the payload while leaving the header unencrypted. Examples of this are
412 JOSE (JSON Object Signing and Encryption) [28], and its lightweight ver-
413 sion COSE (CBOR Object Signing and Encryption) [29]. These encryption
414 mechanisms are also the basis of key exchange protocols such as EDHOC
415 (Ephemeral Diffie-Hellman Over COSE) [30] and application-layer security
416 schemes like OSCORE (Object Security for Constrained RESTful Environ-
417 ments) [31]. Because of their optimisation for constrained environments, this
418 paper focuses on the combined use of COSE, EDHOC and OSCORE as the
419 potential object security solutions for Industry 4.0.

420 Another aspect to be addressed in E2E security is the possibility of par-
421 ties outside the OT network having to access the data generated in it. This
422 data retrieval will occur in the DMZ, as explained in Section 3, while confi-
423 dentiality still has to be preserved. To this end, it would prove useful to have
424 an encryption mechanism that enables multiple users to access the informa-
425 tion without re-encrypting it repeatedly or distributing new keys. This can
426 be accomplished with Functional Encryption [26] —i.e., IBE (Identity-Based
427 Encryption) [32] and ABE (Attribute-Based Encryption) [33]. These ciphers
428 encrypt information according to a set of identities (IBE) or attributes (ABE)
429 that users must possess if they want to decrypt it. ABE can therefore be con-
430 sidered an evolution of IBE, since it provides more flexibility by encrypting
431 data in a more detailed manner. This article will cover ABE since attributes
432 provide a more flexible way of defining who is allowed read encrypted data.

433 Summarising, efficient lightweight communication and encryption proto-
434 cols are required in OT networks. In this context, object encryption combined
435 with lightweight data formats provides a compromise between security and
436 computational cost, and can be integrated into the Application and Data
437 layers of the proposed DiD strategy. Section 5 focuses on this possibility.
438 Meanwhile, Section 6 presents a detailed description of attribute-based en-
439 cryption, which provides role-based access to ciphertexts. This allows them
440 to be shared with different endpoints without the user that encrypts data
441 identifying those endpoints one by one, but guaranteeing data confidential-
442 ity.

443 5. Object Security

444 The aim of object security is the protection of the message itself, provid-
445 ing fine-grain access control of its content. This is achieved using “Secure
446 Objects”, which are information containers comprised of a header, an en-
447 crypted payload and an integrity verification tag [27]. The same message
448 may carry several objects, or different parts of the message can be individu-
449 ally protected. Thanks to this property, object security is an effective way to
450 obtain E2E security through middleboxes, since messages can be encrypted
451 so that middleboxes can only read the required information. Therefore, even
452 if intermediate nodes are compromised, payload confidentiality is not jeop-
453 ardized. The object security method for constrained environments proposed
454 by the IETF Working Group, CoRE, is OSCORE. It uses the CBOR data
455 format, COSE for encryption and EDHOC as the key management protocol.
456 They are explained in the following sections.

457 5.1. CBOR

458 The need for an object data format for constrained devices arose with
459 the presentation of the Object Security Architecture for the IoT (OSCAR)
460 [34]. This architecture had low energy consumption, low latency and ensured
461 security through middleboxes. However, it did not include an object security
462 format suitable for constrained devices, so the architecture’s efficiency was
463 reduced in such scenarios [27]. To solve this, the IETF proposed CBOR
464 [35], a data format optimised for highly constrained environments. It uses a
465 binary type data format, which reduces human-readability, but increases the
466 message transmission and coding/decoding speeds.

467 5.2. COSE

468 COSE [29] was proposed to provide CBOR with security mechanisms,
469 such as the creation and processing of signatures, message authentication
470 codes and encryption. It specifies which signature algorithms shall be ap-
471 plied and how to build, encrypt and decrypt messages. COSE messages are
472 constructed in “layers”, allowing for the sought fine-grain-level approach.
473 The standard offers different encryption and signing possibilities, but when
474 working with OSCORE, it only uses the untagged COSE_Encrypt0 structure.

475 This protocol does not specify the recipients of the message and assumes
476 that they know the key to be used for decryption. Therefore, it should be
477 combined with key management protocols like EDHOC.

478 *5.3. EDHOC*

479 EDHOC is a lightweight key exchange protocol with a small message
480 overhead [30], making it efficient for technologies with duty-cycle or battery
481 limitation. According to the standard, EDHOC also provides the following
482 security features:

- 483 • Mutual authentication with aliveness. This means that the communi-
484 cating parts authenticate each other. This way, both endpoints know
485 they are communicating with whom they intended. It helps reduce
486 impersonation attacks.
- 487 • Perfect Forward Secrecy (PFS). EDHOC achieves this by running an
488 Elliptic Curve Diffie-Hellman (ECDH) key exchange with ephemeral
489 keys. It guarantees that if an attacker gets the keys, it only gets the ones
490 being used in the moment of an attack, and every message exchanged
491 with previous keys continues to be confidential.
- 492 • Identity protection. Passive attackers cannot learn the identity of ei-
493 ther communicating party. Active attackers can only learn about the
494 receiver [36].
- 495 • Crypto Agility, given by COSE. This facilitates changing the cryptog-
496 raphy algorithms, making potential system upgrades faster and easier.
- 497 • Protection against replay attacks. This prevents attackers from re-
498 sending messages that have already been received.
- 499 • Protection against message injection. This prevents an attacker from
500 injecting fake messages into the stream.

501 Although EDHOC does not add requirements to the transport layer it
502 is recommended to implement it in combination with CoAP [37], CoRE's
503 communication protocol for constrained devices. They have also developed
504 a draft with new configuration options to improve CoAP default security,
505 including the prevention of amplification attacks. Its implementation is en-
506 couraged to prevent IIoT devices from being manipulated to launch DDoS
507 attacks. The interested reader is referred to [38] for more details about these
508 enhancements.

509 EDHOC key exchange takes three messages between a Party U (initiator)
 510 and a Party V (responder), after which message exchange between both par-
 511 ties is protected. Each of these three messages is a CBOR sequence protected
 512 by COSE. EDHOC supports various authentication methods—i.e., certifi-
 513 cates, PSK (pre-shared keys) and RPK (raw public keys). The parameters
 514 exchanged between parties will vary between methods, but a simplification
 515 is included in Figure 5.

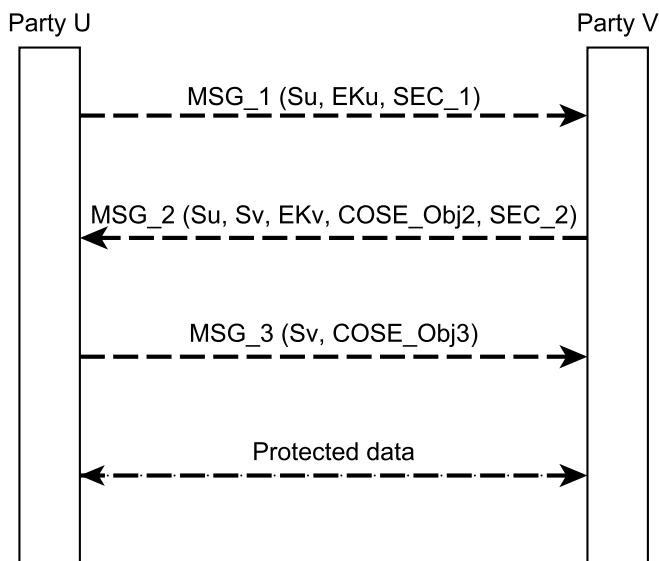


Figure 5: EDHOC negotiation messages.

516 In Figure 5, MSG_1 includes party U’s session key (Su) and ephemeral
 517 key (EKu), and SEC_1. SEC_1 specifies the supported elliptic curves for the
 518 ECDH as well as the supported cipher suites. MSG_2 answers with both
 519 party’s session keys (Su and Sv), V’s ephemeral key (EKv), COSE_Obj2 and
 520 SEC_2. SEC_2 now contains the selected elliptic curves and cipher suites.
 521 Finally, MSG3 contains Party V’s session key and COSE_Obj3. As it is sum-
 522 marised in [39], COSE_Obj2 is used to protect MSG_1 and MSG_2 integrity,
 523 and to authenticate the server. Meanwhile, COSE_Obj3 authenticates the
 524 client and ensures the integrity of the exchanged messages.

525 The security features of EDHOC are in line with the security require-
 526 ments for Industry 4.0 detailed in Section 2. For instance, the protection

527 against both replay and message injection attacks may prevent an attacker
528 from sabotaging the control messages. Moreover, since it provides perfect
529 forward secrecy, EDHOC helps to mitigate pervasive monitoring, preventing
530 an attacker from learning more about the system to prepare a more harmful
531 attack. Finally, the first message exchanged in EDHOC allows verifying that
532 the chosen cipher suite is supported by both communicating parties, which
533 is necessary in the commonly heterogeneous manufacturing environments.

534 5.4. OSCORE

535 OSCORE [31] is CORE’s application-layer security framework for con-
536 strained environments. It uses EDHOC as key exchange protocol and pro-
537 tects messages using COSE. Integrity and confidentiality are provided by
538 the Authenticated Encryption with Associated Data algorithm (AEAD) [40],
539 while authentication and authorisation come from using the Authentication
540 and Authorisation for Constrained Environments (ACE) standard [41].

541 OSCORE also improves COSE’s security by encrypting the method in
542 the original header and placing it in the encrypted payload. A dummy code
543 is then placed in the new header: POST for requests and CHANGED for
544 responses. This prevents attackers from changing a PUT to a DELETE and
545 deleting a resource. Figure 6 shows how OSCORE messages are built upon
546 CoAP messages. Some fields are encrypted, others only integrity protected,
547 and others are left in plaintext (box 2). This information is encapsulated in
548 a COSE message (box 3), which is the content of the ciphertext field of the
549 OSCORE message (box 4). Therefore, the payload is now encrypted, while
550 the header fields remain in plain text and can be processed by middleboxes,
551 if necessary.

552 Apart from providing E2E security even in the presence of middleboxes,
553 OSCORE guarantees most of the industrial security requirements specified
554 in Section 2.2. These include integrity, authentication and authorisation.
555 Moreover, OSCORE is specially designed for constrained networks, making
556 it highly optimised for IIoT nodes. As shown in [42], it has less overhead
557 than CoAP+DTLS, it is faster both in single-hop and multiple-hop scenarios,
558 and it also deals better with retransmissions. Finally, the combined use of
559 OSCORE and EDHOC has a small footprint [30], thanks to the fact that
560 both use CBOR and COSE.

561 The use of these protocols, specially designed for constrained devices,
562 make OSCORE very useful for securing messages between the IIoT nodes
563 constituting an OT network. Furthermore, EDHOC provides the perfect

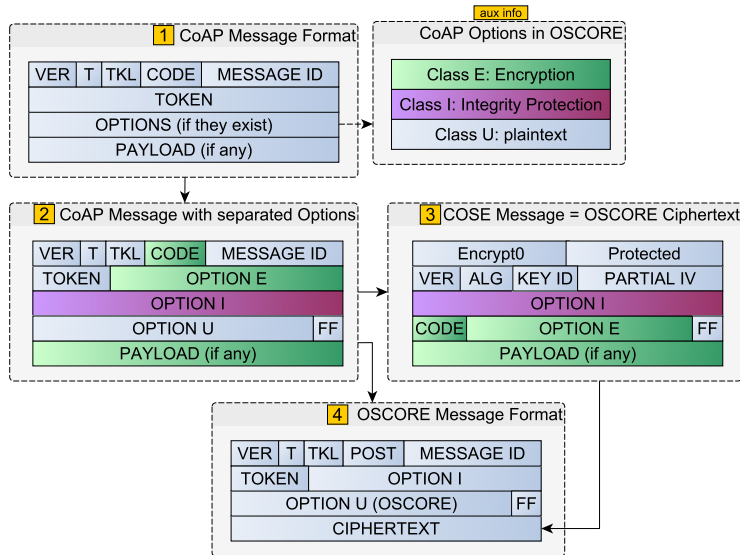


Figure 6: Composition of OSCORE messages.

564 forward security OSCORE cannot provide by itself. In case the keys are
 565 compromised, this property ensures that every encrypted message exchanged
 566 in previous sessions remains protected. Industry 4.0 will also benefit from
 567 OSCORE's header compression and it being mappable to HTTP. The compression
 568 reduces the per-packet overhead, making the transmission of industrial
 569 small data packets faster. The compatibility with HTTP facilitates the
 570 connectivity IIoT nodes need.

571 6. Attribute-Based Encryption

572 As shown in Section 5, OSCORE protects requests and responses using
 573 partially encrypted messages. It also uses CoAP as the communication
 574 protocol, which supports requests to an IP multicast group [43]. However,
 575 protecting group messages with OSCORE [44] entails challenges such as handling,
 576 distributing and updating keys. As a result, the efficiency of OSCORE
 577 is reduced in situations where data needs to be encrypted and distributed
 578 to a group whose members change frequently. ABE can solve this issue by
 579 relating ciphertexts to attributes. In Industry 4.0, it may be applied to confidential
 580 or sensitive information that has to be accessed by parties from
 581 outside the OT network. This can be the case of audit logs [45]: each entry

582 could be encrypted according to an access policy, giving different endpoints
583 particular access rights to the same bulk of data without worrying about key
584 distribution.

585 Because ABE creates ciphertexts according to a set of attributes or roles,
586 senders do not need to know the identity of every recipient. This allows
587 data to be encrypted once and shared with multiple users, simplifying key
588 management in comparison with OSCORE. For instance, in a publisher-
589 subscriber communication model (e.g., MQTT, AMQP or CoAP Pub/Sub
590 [46]) the use of ABE means that the group key does not have to be updated
591 or the information re-encrypted whenever a new node joins the network,
592 improving scalability [47]. This makes ABE a very interesting encryption
593 mechanism for Industry 4.0.

594 In ABE a user with a private key ω may decrypt data encrypted with
595 the public key ω' , if and only if the difference between ω and ω' is minimal
596 [33]. What constitutes these keys depends on whether the chosen approach is
597 Key-Policy ABE (KP-ABE) [45] or Ciphertext-Policy ABE (CP-ABE) [48].
598 In KP-ABE the plaintext is encrypted according to a subset of attributes.
599 Meanwhile, in CP-ABE the plaintext is encrypted according to a policy that
600 dictates which attributes must be fulfilled to decrypt the message. CP-ABE
601 is more interesting for Industry 4.0 applications because it gives the sender
602 of the message full control over who will be capable of decrypting it. This is
603 called implicit authorisation, and it works as follows:

- 604 1. Private keys are associated with an arbitrary number of attributes ex-
605 pressed as strings. For example:
606 - A database in Security Zone A has the attributes: $\{\text{"Zone A"} \wedge$
607 $\text{"Database"}\}$.
608 - A robotic cell in Security Zone A has the attributes: $\{\text{"Zone A"} \wedge$
609 $\text{"Robotic cell"}\}$.
610 - A database in Security Zone B has the attributes: $\{\text{"Zone B"} \wedge$
611 $\text{"Database"}\}$.
- 612 2. The ciphertext specifies an access policy/structure over a defined uni-
613 verse of attributes within the system. The policy is established by the
614 sender. For example, a temperature sensor sends readings with the
615 following access structures:
616 - Temp. 01: $\{\text{"Zone A"} \wedge (\text{"Database"} \vee \text{"Robotic cell"})\}$
617 - Temp. 02: $\{\text{"Zone A"} \wedge \text{"Database"}\}$
- 618 3. The recipient may decrypt the ciphertext if and only if its attributes
619 fulfil the ciphertext's access structure.

620 In this case, the database in Security Zone A is able to decrypt both
621 temperatures, the robotic cell is only able to decrypt the first one, and
622 the database in Security Zone B can decrypt neither.

623 In an Industrial environment, ABE achieves E2E security and provides
624 role-based access control to data. In Industry 4.0, it is becoming more usual
625 for entities outside the OT network to need access to the data generated in it.
626 The privileges of these entities have to be controlled and limited according
627 to their needs. Using ABE over CoAP to encrypt the information provided
628 to these entities ensures that only legitimate endpoints can decrypt it.

629 Finally, integrating ABE in a DiD framework should be straightforward.
630 DiD calls for role-based access whenever possible, and thus the structures to
631 define the access policies should already be in place. Therefore, these trusted
632 entities can also be used to distribute the original attributes of ABE.

633 7. CONCLUSIONS

634 This paper presents an overview of security measures and recommenda-
635 tions for a secure Industry 4.0, where E2E security is usually not guaranteed
636 in the presence of some intermediate elements, such as proxies or gateways.

637 First, best practices to secure Industry 4.0 are identified. They aim to
638 enhance OT network security by adapting and implementing IT security rec-
639 ommendations. These suggestions focus on applying traditional IT security
640 requirements to Industry 4.0. They involve authentication, confidentiality,
641 integrity, availability and non-repudiation. However, most Industry 4.0 en-
642 vironments will require more sophisticated implementations to meet those
643 requirements. For this reason, a Defence-in-Depth approach is suggested.

644 In a DiD strategy, security is divided in layers to address as many attack
645 vectors as possible. These layers can be adapted to company criteria, but they
646 should guarantee the following: restricted access to the network and IIoT
647 devices, the separation of OT and IT networks and the use of secure protocols.
648 Compliance with these requirements should be reviewed periodically and be
649 accompanied by corporate policies that ensure a rapid restoration of the
650 system.

651 The presented DiD layers, along with the technologies considered for
652 them, comply with security specifications. These procedures include imple-
653 menting role-based access control coupled with the principle of least privilege.
654 To segregate IT and OT, the use of NGFW and DMZ has been proposed.

655 It is also suggested to combine these firewalls with IDS and IPS to monitor
656 inbound and outbound traffic while highlighting the importance of avoiding
657 false positives from IPS. Keeping sensitive information confidential is vital
658 in Industry 4.0, so encryption is integrated into the DiD proposal.

659 The proposed solutions are OSCORE and ABE. OSCORE provides E2E
660 security by encrypting the message payload and leaving the header fields in
661 plaintext. Thus, gateways can process messages without breaking their con-
662 fidentiality. OSCORE is concluded to be an appropriate security framework
663 for Industry 4.0 thanks to its header compression, data format and optimised
664 key exchange protocol. Other features that reinforce this conclusion are its
665 capability of working with HTTP, which reinforces IIoT devices' connectivity.

666 Finally, ABE is the encryption proposed to manage third party access to
667 the information contained in the OT network. Since IIoT nodes are highly
668 structured, and changes are rare and predictable, any outsider temporarily
669 accessing the system is considered a vulnerability in the Industry 4.0 security
670 framework. To counter this, we propose to encrypt the data required by these
671 parties with ABE. This allows fine-grained access control to sensitive data
672 and simplifies key management, avoiding having to issue new keys and to re-
673 encrypt messages whenever a new entity accesses the system. Besides, ABE is
674 determined to have easy integration into the DiD environment. The trusted
675 third-party used to define the roles for role-based access can be employed
676 to determine and distribute the attributes and the access policies for the
677 information to be shared.

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